



Implementation of Triple Transposition Vegenere Cipher Algorithm and Cipher Block Chaining for Encoding Text

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Abstract – Security of confidential text data and personal is very important so that the data is not misused by other parties. Many cryptographic algorithms can be used to secure data to keep the originality of confidential data. This study describes the use of triple transposition vegenere cipher algorithm which is combined with cipher block chaining mode for encoding text data with the aim to optimizing randomize of ciphers generated so that the security of confidential of data text is more optimal.

Keywords – Cryptography, Triple Transposition Vegenere Cipher, Cipher Block Chaining, Data, Text

1 INTRODUCTION

Security of confidential or private data text data nowadays is one of the most important things to do because currently the distribution of information of data is often done. Ignoring securing confidential data is very fatally and can disadvantage the owner of the information itself. One effort that can be done is to utilize cryptographic techniques so that data transmitted through communication channels can be safe[1]. Cryptographic techniques can change two-way communication with the process of encryption and decryption[2]. Using cryptographic techniques can increase data attackers' complexity to find out the original meaning of the data easily.

Data randomization effectively will add complexity to finding the original patterns and meanings of confidential data. This research describes combining the triple transposition vegenere cipher algorithm to perform text encoding and cipher block chaining operation modes used to optimize randomization of encoded text data. Secret text that is secured is encoded based on the triple transposition vegenere cipher algorithm three times the transposition and substitution process, then the resulting cipher binaries will be scrambled again based on the cipher block chaining operation mode.

Implementation of the triple transposition vegenere cipher algorithm and the cipher block chaining operation mode to secure confidential data can generate the more random ciphertext of secret data and can increase the hassle of finding the original pattern and meaning of the text data.

2 THEORY

2.1 Cryptography

Study of mathematical techniques covering data security aspects is generally referred to as cryptography[3][4]. Confusion and diffusion in reconstructing the resulting cipher is the main achievement of cryptographic techniques. Objectives that must be achieved in the application of cryptographic techniques are confidentiality, data integrity, authentication and non-repudiation[5][6].

Cryptographic techniques will encrypt the original data to be a cipher called the encryption process and the decryption process to restore the cipher to be original data. The process of encryption and decryption requires a structured algorithm and the key that used to encryption and decryption processes.

2.2 Triple Transposition Vegenere Cipher Algorithm

Triple Transposition Vegenere Cipher (TTVC) algorithm is the development of the vegenere cipher algorithm. This algorithm will encryption and decryption a data by doing two processes is transposition and substitution. this process is doing three times. Transposition process is doing based on columnar transposition and substitution is doing based on vegenere cipher. Each process is used a different key, meaning that there are three keys for the substitution process and three keys for transposition processes[7].

Workflow of Triple transposition vegenere cipher in the encryption and decryption process can be presented in Figure 1 below.

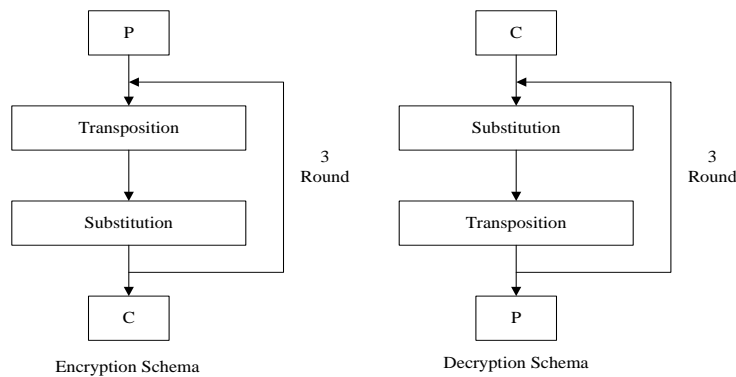


Figure 1. Encryption and Decryption Scheme Based on Triple Transposition Vegenere Cipher
Source : M. Linda Caroline, 2011[8]

Based on Figure 1 above, it can be seen that the original data encryption process is done alternately starting from the transposition process based on columnar transposition (T) and substitution (S) is done based on the vegenere algorithm. This process is carried out in three rounds with different keys for each process. The decryption process is the opposite of the encryption process, the key used remains the same as the key of an encryption process. The formulations for the encryption and decryption process[8] are :

$$\text{Encryption formulation : } C = (S3(T3(S2(T2(S1(T1(P))))))) \dots\dots\dots(1)$$

$$\text{Decryption formulation : } P = (T1(S1(T2(S2(T3(S3(C)))))) \dots\dots\dots(2)$$

Formulation that we use for substitution is vegenere algorithm formulation are :

$$\text{Encryption formulation } = C_i = (P_i + K_i) \text{ mod } 256 \dots\dots\dots(3)$$

$$\text{Decryption formulation } = P_i = (C_i + K_i) \text{ mod } 256 \dots\dots\dots(4)$$

In this case uses 256 as modulus value.

2.3 Cipher Block Chaining Mode

Cipher Block Chaining (CBC) is one of the operating modes commonly used in modern cryptographic algorithms. This operating mode is very superior in randomization binary of data that is loaded in blocks. The cipher block chaining operation mode will perform encryption and decryption process that interdependent between the binary blocks of the original binary data. In this case, the previous binary-block binary will be feedbacked in the process of encryption or decryption of the current block, so the block cipher is generated entirely depends on all blocks biner of plaintext[5][9]. Scheme of encryption and decryption process based on cipher block chaining mode are presented in figure 2 below.

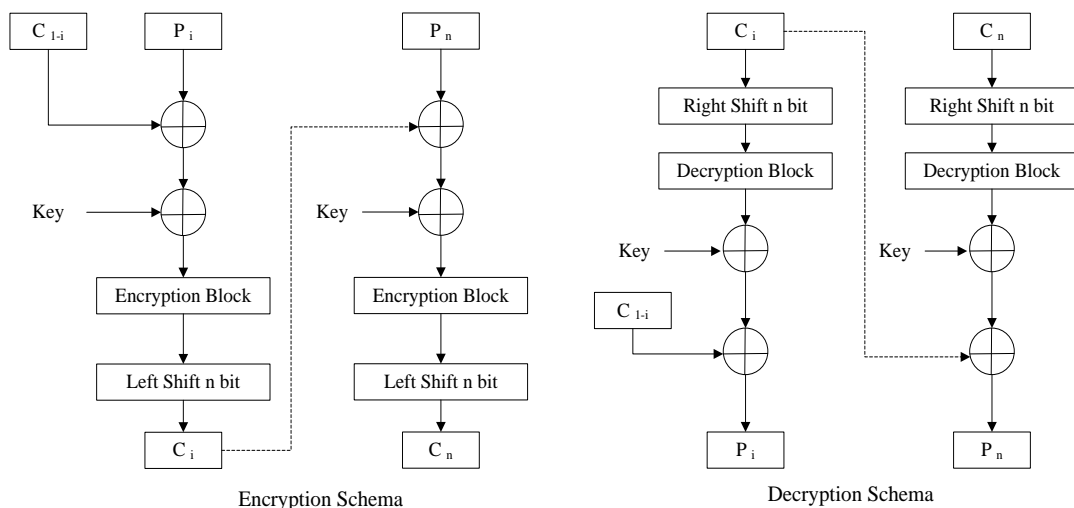


Figure 2. Encryption and Decryption Scheme Based on Cipher Block Chaining
Source : T. Zebua, 2015[5]



Based on the above scheme, it is known that block of plaintext binary and initial cipher (C1-i) block will be processed with XOR. Variable of C1-i in CBC is often called C0 or initial cipher (IV) or initial cipher. Initial cipher and the number of bits is determined by the user as an initialization vector. The number of C0 bits and keys must be same as the length of a plain bit block or cipher. The number of shift bits can be determined by the user included the position of shifted (left to right or right to left), number of bits that shifted as default in CBC is 1 bit.

Displacement of bit in decryption process is reversed position of encryption and the number of bits is same[10]. Based on this description, it can be formulated that:

Encryption Formulation : $C_i = \text{Shift 1 bit to left} ((P_i \oplus C_i) \oplus \text{Key}) \dots\dots\dots(5)$

Decryption Formulation : $P_i = (\text{Shift 1 bit to Right}(C_i) \oplus \text{Key}) \oplus C_i \dots\dots\dots(6)$

3 RESULT AND DISCUSSION

Results of the implementation triple transposition vegenere cipher algorithm that combination with the cipher block chaining mode to secure text data are presented in the figure 3 below.

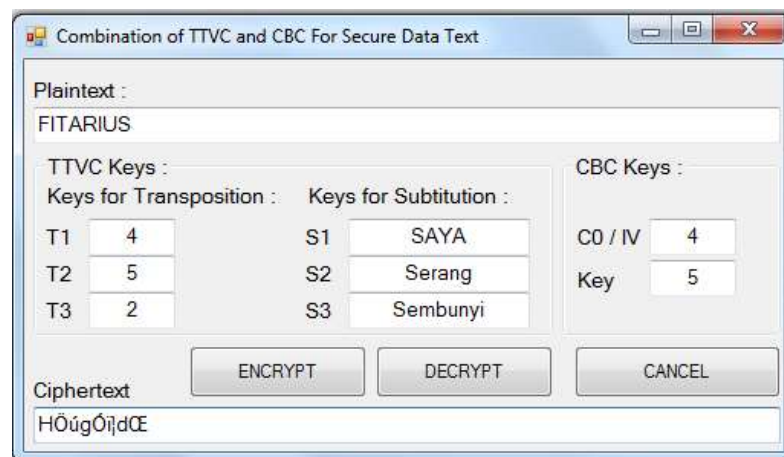


Figure 3. Result of Implementation Combination TTVC algorithm and CBC Operation Mode

Combination process of Triple Transposition Vegenere Cipher (TTVC) and Cipher Block Chaining (CBC) mode for secure the text in this study begins by encrypt the text data based on the triple transposition vegenere cipher algorithm, then binary of cipher that resulted will be randomized based on the cipher block chaining mode. The encryption process can be presented in the diagram below.

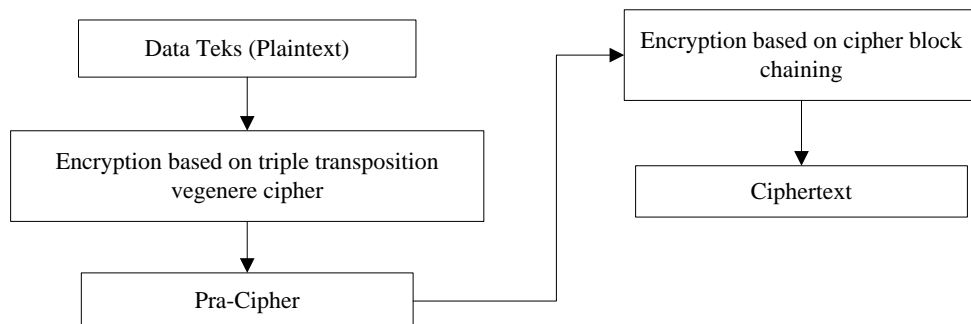


Figure 4. Ecrption Process Scheme Base on Combination of TTVC Algorithm and CBC Mode

Decryption process of ciphertext to be plaintext begins with the decryption process of ciphertext based on the cipher block chaining, then the plaintext obtained from the cipher block chaining decryption process will be decrypted again based on the triple transposition vegenere cipher process. Decryption process diagram is shown in figure 5.

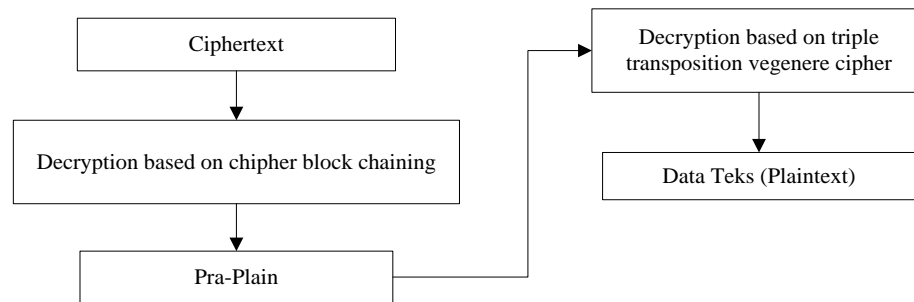


Figure 5. Decryption Process Scheme Based on TTVC Algorithm and CBC Mode

Text that used as samples in this research is FITARIUS, the key used in the encryption and decryption process based on the triple transposition vegenerer cipher algorithm is:

The key for columnar transposition is T1 = 4, T2 = 5, T3 = 2 and keys for vegenerer cipher is S1 = SAYA, S2 = Serang, S3 = Sembunyi

Based on figure 4 above, the encryption process begins based on the triple transposition vegenerer cipher algorithm, then the pre-cipher obtained is scrambled again based cipher block chaining mode.

1. Encryption process based on triple transposition vegenerer cipher (equation 1)

a. First Transposition

Value of key for each transposition (T1, T2, T3) is used as the number of transposition table columns that are formed.

F	I	T	A
R	I	U	S

Cipher retrieval starts in each column, so the first transposition (T1) cipher is generated FRIITUAS.

b. First Substitution

Cipher of first transposition (T1) will be substituted based on the vegenerer cipher algorithm and the key used is the S1 key (follow equation 1), so that the S1 cipher is generated "ç Š§- Š"

Cipher S1 will use as plaintext on second transposition process

c. Second Transposition

TM	“	ç	Š	§
-	Š	”	X	X

Character of X in the 5th and 6th columns in the last row are additional characters to fill in the blank fields. The second transposition cipher is "çŠ§X§X". Cipher T2 is used as plaintext on the second substitution process.

d. Second Substitution

Second substitution using key S2 and cipher T2 is used as plaintext in this process (follow equation 1). Cipher S2 is "Ÿ½¹" and this cipher is used as plaintext in third transposition process.

e. Third Transposition

Ÿ	½
¹	½
½	¹
½	¹
¹	½
½	¹

Cipher T3 is "Ÿ½¹" and is used as plain for third substitution process.

f. Third Substitution

Cipher T3 will be substituted based on vegenerer cipher algorithm with S3 key (follow equation 1), so the cipher (cipher of S3) generated is "Žitd-". This cipher is the final result of encryption process based on triple transposition vegenerer cipher. Binners of this cipher will be random based cipher block chaining mode.



2. Encryption process based on cipher block chaining (follow equation 5)
The number of group bits is 40 bits and $C0 / IV = \text{STMIK}$, Key = HUMEN, Shift 1 bit from left to right.
The cipher generated based triple transposition vegenere cipher encryption process is ?j}?Žitd[-
Biners of this cipher characters is :
001111101101010011111010011111100011100110100101110100011001000001000000011110
Separate these bits into two groups with the number of bits each group is 40 bits and begin from left to right

P1 = 00111110110101001111101001111110001110
P2 = 0110100101110100011001000001000000011110

Encrypt P1, P2 for resulted C1, C2 based on equation 5, then resulted:

C1 = 0100100011010110111110100110011100010110
C2 = 1101001111101111101001100110010010001100

Combine bits of C1 and C2, then split into 8 bits each group and then convert to be a characters:

010010001101011011111010011001100010110100111110111101001100110010010001100

So, the cipher of text that resulted is **HÖúgÓr,dE**

3. Decryption Process
Decryption process begins based on the cipher block chaining decryption (use equation 6). The plaintext obtained from this decryption process will be used as a cipher in the decryption process based on the triple transposition vegenere cipher (see figure 5 above).
- a. Decryption process based on cipher block chaining mode
Value of C0, keys and the shift value used in the decryption process are the same as the values used in the encryption process. It's just that the bit position shifted in the decryption process is the bit in the right position (because bit displacement position on encryption process is from left to right), the number of shifts is the same. The initial stage is to convert the ciphertext into binary, then group the binary according to the number of bits agreed upon in the encryption process. Follow the equation 6 to do the decryption process. Plaintext that resulted will be converted to characters and resulted characters ?j}?Žitd[- (this is the result of encryption process based on triple transposition vegenere cipher).
- b. Decryption process based on triple vegenere cipher
Ciphertext that used in the decryption process based on the triple transposition vegenere cipher is the plaintext that is generated from the decryption process of cipher block chaining. The decryption process is carried out based on equation 2. Value of key that used is the same value as the encryption key value. Based on equation 2, plaintext is resulted is same as before, that is **FITARIUS**.

4 CONCLUSION

The results of analysis of the previous chapters, it can be concluded, where the conclusions are likely to be useful for the readers. Thus, the writing of this article can be more useful. The conclusions are as follows:

1. Combination of triple transposition vegenere cipher algorithm with cipher block chaining mode, able to increase randomness to encoded text of data, so as to be able to hide the patterns and original meaning of data text.
2. Randomizing binary-binary of confidential data based on cipher block chaining mode able to realize confusion and diffusion for the attackers to know the original meaning of data.
3. The combination of triple transposition vegenere cipher with cipher block chaining mode is very well if developed and used to improve the optimization of classic cryptographic algorithms.

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