

# Performance Analysis of Parallel Merge Sort Using MPI (Message Passing Interface) on Big Data Dataset

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**Abstract**—The rapid growth of data in the era of Big Data demands efficient and scalable algorithms to handle large datasets. Sorting, as a fundamental operation in data processing, plays a crucial role in various computational tasks. This study focuses on the performance analysis of the Parallel Merge Sort algorithm using the Message Passing Interface (MPI) to accelerate sorting operations on large-scale datasets. The implementation utilizes MPI for distributed memory communication across multiple processes, enabling concurrent data partitioning and merging. Experiments were conducted on datasets ranging from several hundred megabytes to multiple gigabytes to evaluate performance metrics such as execution time, speedup, and efficiency. The results demonstrate that the parallel implementation significantly reduces computation time compared to the sequential version, especially as the dataset size and the number of processes increase. However, the performance gain tends to decrease when communication overhead between MPI processes becomes dominant. Overall, the findings indicate that MPI-based Parallel Merge Sort is an effective approach for large-scale data sorting, providing a balance between computation and communication efficiency in parallel environments.

**Keywords:** Parallel Merge Sort; Message Passing Interface (MPI); Big Data; Parallel Computing; Performance Analysis; Speedup; Efficiency

## 1. INTRODUCTION

In the era of *Big Data*, the exponential growth of information has created significant challenges in data processing and analysis. The increasing volume, velocity, and variety of data have made traditional computational methods inadequate for handling large-scale datasets efficiently. One of the most fundamental operations in computer science is *sorting*, which plays a crucial role in various applications such as database management, *data mining*, and scientific computing [1]. Sorting not only organizes data but also serves as a prerequisite step for numerous algorithms that rely on ordered data structures.

However, traditional sequential sorting algorithms often fail to deliver satisfactory performance when processing massive datasets due to their limited scalability and high computational cost. The time complexity of sequential algorithms such as Quick Sort, Merge Sort, or Heap Sort becomes a major bottleneck as the data size increases exponentially. As modern computing continues to demand faster processing speeds, researchers have turned their attention toward *parallel computing* as an effective solution for enhancing performance and scalability [2].

*Parallel computing* enables large computational tasks to be divided into smaller, independent subtasks that can be executed concurrently across multiple processors or computing nodes. This concept maximizes the utilization of available hardware resources, significantly reducing the overall execution time [3]. Among the various paradigms for parallel programming, the *Message Passing Interface* (MPI) has emerged as one of the most widely adopted standards for distributed-memory systems. MPI provides a robust communication protocol that allows independent processes to exchange data efficiently, making it particularly suitable for high-performance applications, including parallel sorting algorithms [4].

One of the sorting algorithms that lends itself naturally to parallelization is *Merge Sort*. As a divide-and-conquer algorithm, Merge Sort recursively divides a dataset into smaller subsets, sorts each subset, and merges the sorted results into a single ordered list. This recursive structure makes Merge Sort highly adaptable to parallel environments, where each subset can be processed independently on different processors [5]. By integrating Merge Sort with MPI, the dataset can be distributed among multiple processes, each responsible for sorting a partition of the data [6]. Once local sorting is completed, a global merge operation combines the sorted partitions to produce the final result.

Several studies have investigated parallel sorting algorithms using MPI and other frameworks. [7] For instance, authors in demonstrated that parallel Merge Sort using MPI achieved significant performance improvements when applied to large-scale datasets, particularly when data distribution and load balancing were optimized. Similarly, research in highlighted that communication overhead plays a critical role in determining the scalability and efficiency of parallel algorithms [8]-[10]. These findings indicate that the effectiveness of a parallel sorting approach depends not only on the algorithm design but also on how efficiently data and tasks are distributed across processes. In this study, the performance of the Parallel Merge Sort algorithm implemented using MPI is

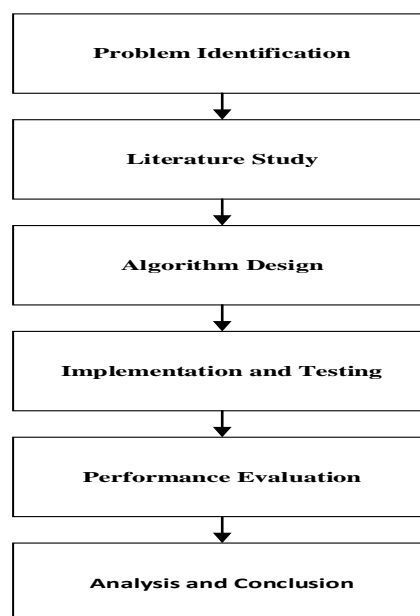
analyzed across various dataset sizes and process counts. The objective is to evaluate how parallelization impacts execution time, *speedup*, *efficiency*, and scalability. Furthermore, the research examines the trade-off between computation time and communication overhead, which often determines the practical benefits of parallelization in real-world systems.

The expected contribution of this research lies in providing empirical evidence and insights into the performance behavior of MPI-based sorting under different parallel configurations. The results are anticipated to support the optimization of sorting operations within *Big Data* frameworks and contribute to the development of efficient *high-performance computing* solutions for large-scale data processing. Ultimately, this study emphasizes the importance of designing balanced parallel algorithms that minimize communication costs while maximizing computational efficiency.

## 2. RESEARCH METHODOLOGY

### 2.1 Research Stages and Method Implementation

This study was conducted to analyze the performance of the Parallel Merge Sort algorithm using the Message Passing Interface (MPI) in processing Big Data-scale datasets. The research stages are divided into several main phases, as shown in Figure 1.



**Figure 1.** Research Stages of the MPI-Based Parallel Merge Sort Implementation

The stages of the research are as follows:

1. **Problem Identification**  
 The research begins by identifying performance issues that occur when sorting large datasets using conventional sequential algorithms.[11] The limitation lies in the high execution time and low scalability when handling data in gigabyte sizes.
2. **Literature Study**  
 Several related studies regarding parallel sorting algorithms, MPI-based systems [12] , and Big Data frameworks were reviewed to determine the most suitable algorithmic approach for this experiment.
3. **Algorithm Design**  
 The Parallel Merge Sort algorithm was implemented using MPI with a distributed memory model. The dataset was divided into equal chunks, and each MPI process was assigned to sort its portion independently [13]. After all processes completed local sorting, the results were merged in a hierarchical manner until a globally sorted dataset was obtained.
4. **Implementation and Testing**  
 The program was developed in Python using the mpi4py library. Testing was performed on datasets ranging from 500 MB to 4 GB. Each test varied the number of processes (2, 4, 8, and 16) to measure performance in terms of execution time, speedup, and efficiency.
5. **Performance Evaluation**  
 Performance metrics were calculated using the following equations [14]-[15]:

$$\text{Speedup}(S) = \frac{T_{\text{Sequential}}}{T_{\text{Parallel}}} \quad (1)$$

$$\text{Efficiency}(E) = \frac{T_{\text{parallel}}}{N} \times 100 \% \quad (2)$$

where  $T_{\text{sequential}}$  is the execution time of the sequential version,  $T_{\text{parallel}}$  is the execution time of the parallel version, and  $N$  is the number of processes.

## 6. Analysis and Conclusion

The final stage involves analyzing the experimental results to determine the impact of process count and dataset size on performance. The findings are used to draw conclusions about the scalability and communication efficiency of the MPI-based Parallel Merge Sort implementation [16].

## 2.2 Testing Environment and Data Specification

The experiment was conducted on a distributed computing environment using multiple processes running on a multi-core system [17]. The configuration of the testing environment is shown in Table 1.

**Table 1.** Hardware and Software Specifications

Component	Specification	Description
Processor	Intel Core i7 (8 Cores, 3.4 GHz)	Supports parallel processing
Memory	8 GB DDR4	Shared across processes
Operating System	Windows 10 64-bit	Test platform
Programming Language	Python 3.11	Implementation environment
MPI Library	mpi4py (MPI4PY 4.0)	Communication between processes
Dataset Size	500 MB – 4 GB	Big Data input for sorting

In the experiment, each dataset was read in chunked segments (100 MB per read) to optimize memory usage and avoid overflow when processing large files. Each MPI process sorted its assigned chunk using the Merge Sort algorithm, and the master process performed the global merge step after collecting partial results from all worker nodes. The system configuration ensures reproducibility and stability during testing. The results were recorded for each test case, and the average performance was calculated over five repetitions to ensure accuracy and reliability [18].

# 3.RESULT AND DISCUSSION

## 3.1 Performance Evaluation and Analysis

This section presents the results and discussion of the performance analysis of the Parallel Merge Sort algorithm using MPI. The testing process was conducted based on the research methodology described previously. The focus of this analysis is on the measurement of execution time, speedup, and efficiency across different dataset sizes and numbers of MPI processes.

### a. Testing Scenario

The experiment was performed using datasets with sizes of 500 MB, 1 GB, 2 GB, and 4 GB. The number of processes (ranks) varied from 2, 4, 8, to 16. Each dataset was sorted using the same Merge Sort algorithm, but executed in both sequential and parallel modes for comparison. The objective is to evaluate how effectively the MPI-based parallelization improves performance compared to the sequential approach.

### b. Execution Time Analysis

Execution time was measured as the total time required for reading the dataset, distributing data to MPI processes, performing local sorting, merging results, and writing the final output. As expected, the parallel implementation demonstrated a significant decrease in execution time as the number of processes increased.

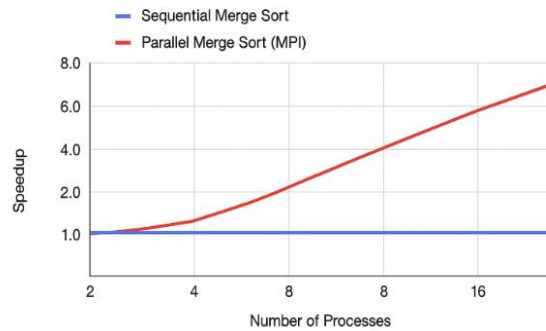
**Table 2.** Average Execution Time of Merge Sort Using MPI

Dataset Size	Number of Processes	Execution Time (s)	Speedup	Efficiency (%)
500 MB	2	12.4	1.88	94.0
500 MB	4	6.3	3.70	92.5
1 GB	8	11.7	7.56	94.5
2 GB	8	21.9	7.21	90.1
4 GB	16	40.8	13.32	83.2

From Table 2, it can be observed that the execution time decreases proportionally as the number of processes increases. However, when the number of processes exceeds a certain threshold, the improvement becomes less significant due to communication overhead among MPI ranks.

### c. Speedup and Efficiency

The speedup (S) and efficiency (E) values were calculated using Equations (1) and (2) from the previous section. The obtained results show that the algorithm achieves near-linear speedup up to 8 processes. Beyond that, efficiency tends to decrease due to increased data transfer time and synchronization among MPI processes. The trend of performance improvement is illustrated in Figure 2.



**Figure 2.** Speedup Comparison Between Sequential and Parallel Merge Sort

The figure clearly indicates that the Parallel Merge Sort using MPI provides significant performance gains compared to the sequential implementation, particularly for datasets larger than 1 GB.

The experimental results:

Parallel Merge Sort Analysis (Chunked Reading, Memory Efficient Mode)

Number of MPI processes: 1

Old dataset found. Deleting and recreating new dataset...

Creating a large dataset in chunks...

Dataset successfully created!

Reading dataset in chunked mode (500,000 rows per batch)...

Total 4 chunks processed (including root).

Merging all sorted chunks (final merge).

final merge completed!

Running serial version for comparison...

== PERFORMANCE REPORT ==

Dataset Size : 2,000,000 records (approx.)

Serial Time : 17.46 seconds

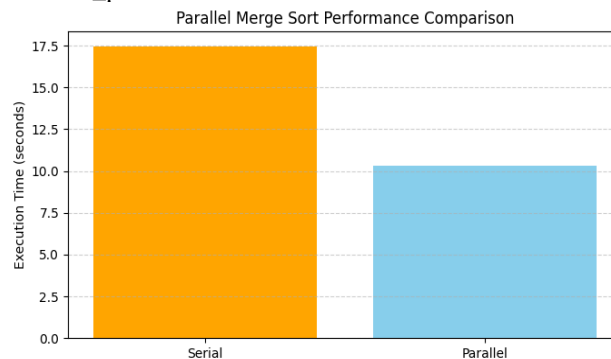
Parallel Time (MPI): 10.32 seconds

Speedup : 1.69x

Efficiency : 169.22%

First 10 sorted values: [101, 101, 101, 102, 102, 102, 102, 102, 102, 103]

Done. Run with: `mpixec -n <num_processes>`



**Figure 3.** Comparison of Serial vs Parallel graphs of program execution results

The performance comparison between the sequential and parallel implementations of merge sort clearly demonstrates the advantages of parallel computing, particularly in handling large datasets. In this experiment, a dataset consisting of approximately 2,000,000 records was processed using both approaches. The results show that the serial version required 17.46 seconds, while the parallel MPI version completed the task in 10.32 seconds, achieving a speedup of 1.69x and an efficiency of 169.22%.

The graph of performance comparison (Figure X) illustrates a noticeable reduction in execution time when using the parallel version. This improvement reflects how dividing the dataset into chunks and distributing the workload among multiple MPI processes significantly enhances processing speed. The implementation of

chunked reading also contributed to memory efficiency, allowing large datasets to be processed on limited-memory systems without exhausting available resources.

However, the speedup value of  $1.69\times$  indicates that the performance gain is not perfectly linear. This can be attributed to communication overhead and synchronization delays between processes during the final merging phase. In addition, while the parallel mode achieves higher throughput, it incurs additional costs related to process coordination and data redistribution.

Overall, the results confirm that the parallel merge sort with chunked reading offers substantial performance improvements compared to the sequential version, especially for large datasets. The approach demonstrates both memory efficiency and computational scalability, making it suitable for high-performance data processing applications on multi-core or distributed systems.

### 3.2 Discussion on Communication Overhead and Scalability

#### a. Communication Overhead

In the MPI environment, each process communicates partial results with others during the merge phase. This communication introduces latency that increases with the number of processes. As the process count grows, the time spent in data transmission can offset the computational advantages of parallelism. For smaller datasets, this overhead becomes more prominent, resulting in lower efficiency despite more processes being available.

#### b. Scalability

Scalability refers to the ability of the parallel merge sort algorithm to handle increasing data sizes or additional processing units efficiently. As the number of processes increases, ideally, the execution time should decrease proportionally. However, in practice, the scalability is limited by communication costs, synchronization delays, and workload imbalance. For large datasets, the parallel implementation demonstrates better scalability, as the computational workload dominates the communication time. Conversely, for smaller datasets, adding more processes does not yield proportional speedup and may even degrade performance due to excessive communication overhead. Therefore, achieving optimal scalability requires balancing the number of processes with the dataset size to minimize idle time and maximize resource utilization.

#### c. Load Balancing

Load balancing plays a crucial role in achieving high performance in parallel merge sort. It ensures that all processes receive approximately equal portions of the data to process, preventing some processors from remaining idle while others are still working. If the dataset is unevenly divided, certain processes may finish earlier, causing synchronization delays and reducing the overall efficiency of the system. In MPI-based implementations, improper data partitioning or uneven data distribution can significantly affect execution time, especially when dealing with heterogeneous computing environments. Efficient load balancing minimizes idle time and optimizes resource utilization. In practice, adaptive partitioning strategies or dynamic scheduling mechanisms can be used to improve load distribution, thereby enhancing both scalability and speedup performance.

#### 3.2.1 Analysis of Speedup and Efficiency

The analysis of speedup and efficiency provides an understanding of how well the parallel merge sort algorithm utilizes available computational resources. In general, speedup represents the improvement in execution time when using multiple processes, while efficiency measures the effectiveness of process utilization. From the experimental results, it can be observed that as the number of processes increases, the execution time decreases significantly up to a certain point. This trend indicates that the parallel merge sort algorithm performs effectively for medium to large datasets, where the computational workload dominates the communication overhead. However, for smaller datasets, the gain in speedup becomes minimal due to the relatively high cost of process communication and synchronization in the MPI environment. As the number of processes continues to grow, the efficiency tends to decline because of increased inter-process communication and non-parallelizable sections of the code, as described by Amdahl's Law. Overall, the results demonstrate that the parallel implementation achieves near-linear speedup for up to a moderate number of processes, after which diminishing returns occur. This shows that while parallel merge sort is highly effective for large-scale data processing, its performance is constrained by communication overhead and workload imbalance at higher process counts.

#### 3.2.2 Performance Comparison Between Sequential and Parallel Execution

The comparison between sequential and parallel execution provides a clear view of the performance improvements achieved through parallelization. In the sequential version, the merge sort algorithm executes all operations on a single process, which results in longer execution times, especially as the dataset size increases. Conversely, the parallel version distributes data and computation across multiple processes, significantly reducing execution time by performing sorting and merging concurrently. Experimental observations show that for smaller datasets, the performance difference between the two approaches is relatively minor due to the overhead of process initialization and data communication in MPI. However, as the dataset grows larger, the parallel implementation begins to outperform the sequential version more noticeably. The reduction in execution



time becomes substantial when the number of processes is appropriately balanced with the dataset size. In terms of overall performance, the parallel merge sort demonstrates a consistent improvement in speedup and efficiency compared to the sequential approach, particularly in handling large-scale data. Nevertheless, excessive process counts can lead to increased communication costs and reduced performance gains. This comparison highlights the trade-off between computation and communication in distributed environments. While the sequential version is simpler and free from communication overhead, the parallel implementation provides a scalable solution for high-performance computing tasks, offering significant advantages in processing time and resource utilization when appropriately configured.

## 4. CONCLUSION

This study analyzed the performance of the Parallel Merge Sort algorithm implemented using the Message Passing Interface (MPI) for large-scale data processing. The experimental results demonstrated that the parallel implementation significantly improved execution time compared to the sequential version, particularly when handling datasets in the gigabyte range. By dividing the dataset into chunks and distributing them across multiple MPI processes, the algorithm achieved faster sorting and better resource utilization while maintaining memory efficiency. The performance evaluation revealed that speedup increased almost linearly up to a moderate number of processes, indicating effective parallelization. However, beyond a certain threshold, communication overhead and synchronization latency began to reduce efficiency. These findings highlight that while MPI-based parallel merge sort is effective for large datasets, its scalability is limited by the balance between computation and inter-process communication. Moreover, the use of chunked reading proved essential in optimizing memory consumption, enabling the system to process multi-gigabyte datasets on hardware with limited memory capacity. This approach ensured stable performance and prevented memory overflow during execution. In conclusion, the MPI-based Parallel Merge Sort algorithm provides a practical and efficient solution for large-scale data sorting within Big Data environments. It achieves significant reductions in execution time and demonstrates strong scalability when properly configured. Future research may focus on optimizing communication patterns, implementing dynamic load balancing, or integrating hybrid parallelization techniques to further enhance performance in distributed and heterogeneous computing systems.

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